

Random Number Generators

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True RNG $\rightarrow$ slow. Expensive

Pseudo RNG $\rightarrow$ Deterministic Algorithms

$\hookrightarrow$ Deterministic sequence of numbers that approximately random

Goal: Create uniform distribution between 0 and 1.

$$U[0, 1)$$

$$\text{reality: } \{0, \dots, m-1\}$$

$$x \sim \{0, \dots, m-1\}$$

$$u = \frac{x}{m-1} \sim U[0, 1]$$

$$= \frac{x}{m} \sim U[0, 1]$$

Definition

Pseudo Random Number Generator

struct (S, m, f, U, g)

S = finite set of states ($m = |S|$)

U = a distribution on S for the initial state (seed)

f = function $S \rightarrow S$

f : $S \rightarrow S$, state transition function

U = output space, $[0, 1)$, or perhaps $\{0, \dots, m-1\}$

g : $S \rightarrow U$, output function

How do we use PRNG?

Select $x_0 \sim U$, $u_0 = g(x_0)$

1. For $n \geq 1$

$$x_n = f(x_{n-1})$$

$$u_n = g(x_n)$$

PRNG are finite state machines $\rightarrow$ they WILL repeat in finite time

$$\text{period} = m$$

p.s.

$$\forall n \geq k \quad S_{p+n} = S_n$$

2. Stop at $\vdash$ goto step

$u_0, u_1, \dots, u_n$ are PR numbers

BAYESIAN SIGNAL PROCESSING

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Three families of PRNGs

1. Linear-congruential generators (LCG)
2. Linear-Feedback shift registers (LFSR)
3. Cellular Automata (CA)

LCG

Works for a long time.

$$x_{n+1} = (a x_n + b) \bmod m$$

$u_n = \frac{x_n}{m}$ $\uparrow$ multiply modulus $\rightarrow$ increment

DEFINITION
MODULO FUNCTIONS
$r = x \bmod m$
$x = km + r$
$\uparrow$ Remainder

~~Theorem: Hull-Dobell Theorem~~

LCG with a, c, m, x_0 has period m

iff: 1. c is relatively prime to m
"coprime"

2. $b = a - 1$ is a multiple of p
for every prime p in m .

3. b is a multiple of 4
if m is a multiple of 4.

MCG's exist. $\rightarrow$ multiplicative / mixed congruential generators

$c=0$

$$x_{n+1} = a x_n \bmod m \rightarrow \text{max period } m-1$$

$\rightarrow$ RANDU from IBM

LFSR

- Fast!
- XOR's
- shift gates

Steps!

1. $x = x \text{ XOR } (x \ll 13)$
2. $x = x \text{ XOR } (x \gg 17)$
3. $x = x \text{ XOR } (x \ll 15)$

MERSENNE TWISTER

Uses rand no: in bitshift $\rightarrow$ Python RAND!
Why XOR and shifts?

- XOR'ing different bit patterns you get complex mixing
- creates an avalanche effect

Adv.

- very fast
- tiny state
- simple to implement
- good enough

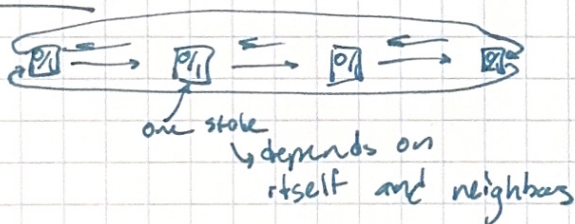
Dis.

- Fails some rigorous tests (Bigcrush, TestU01)
- Not suitable for crypto

PCG: Permuted Congruential Generator
(convolutional?)

combines LCG step plus LFSR step

CA: Cellular Automata



Rule 30

111	110	101	100	011	010	001	000
0	0	0	1	1	1	1	0

$$u = \sum_{i=1}^{n-1} b_i \cdot 2^{-i}$$

- Good for multiple threads
- Parallel process.
- chaotic

